

E3SM All Hands: NGD: Efficient Communication Pattern for Interpolation Semi-Lagrangian Tracer Transport

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Outline





SL MPI

Outline





2 SL MPI

COMPOSE: Semi-Lagrangian Tracer Transport for E3SM



- Problem: Tracer transport is expensive.
- Solution: Semi-Lagrangian (SL) transport. Long time steps; less communication.
- Problem: Transport requires property preservation, and SL makes that harder.
 - ► (mass conservation, shape preservation, tracer consistency, linear correlation preservation)
- Solution: CEDR: Property preservation in exactly 1 all-to-all reduction equivalent.¹
- Opportunity: CEDR enables using the fastest, lowest-communication SL method there is: Interpolation SL (ISL) with compact stencil.
- Problem: ISL based on high-order compact stencil (GLL element, $n_p \ge 4$) is unstable.
- Solution: Stabilized ISL.
- Problem: HOMME's deterministic halo exchange is suboptimal for SL.
- Solution: ISL-specific optimal communication pattern.
 - ► This project.

¹A. M. Bradley, P. A. Bosler, O. Guba, M. A. Taylor, G. A. Barnett, *Communication-efficient property preservation in tracer transport*, to appear in SIAM J. Sci. Comp. Software: github.com/E3SM-Project/COMPOSE



Strong scaling HOMME: Status for 40 tracers



- preqx dycore is >2.1× faster on KNL at 1350 nodes (strong-scaling limit).
- preqx dycore is >3.2× faster on Edison at 3600 nodes (strong-scaling limit).

Outline







Concepts



- Communicate only with ranks holding relevant data.
- For Interpolation SL (ISL), send requests for interpolation.
 - ► Send/receive departure point data to make requests for interpolation.
 - Send/receive interpolation and bounds data, fulfilling these requests.
 - ► In both rounds, overlap communication and computation.



SL MPI speedup



- Prescribed velocity field driver on Mutrino KNL, HSW; time step requires only 1-halo.
- Vary number of elements per rank, number of tracers.
- Measure end-to-end time using original MPI and new SL MPI communication patterns and report speedup.



2-halo overhead



- Prescribed velocity field driver on Mutrino KNL; time step requires only 1-halo.
- Vary number of elements per rank, number of tracers.
- Measure extra time taken when 2-halo is active and report percent overhead.
- Except for this overhead², speedup of 2-halo is exactly the factor time step increase.



²still possibly to do: more accurate velocity, which will also have overhead

2-halo convergence (correctness)





Fidelity study





- Nondivergent flow test case
- "HOMME tuned" data are from Guba et al, *Optimization-based limiters for the spectral element method*, JCP 2014.
- SL MPI is used with 2-halo capability.



- Algorithm: Communicate and compute once per node (currently redundant for edge nodes).
- Algorithm: Expose all available parallelism for GPU implementation; rework metadata and metadata setup as necessary.
- Software: Rewrite several target-cell, local-mesh data structures so bulk data are shared.
 - Bulk data memory footprint independent of halo size.



Thanks!

Semi-structured messages



- Set of departure points in remote elements is nondeterministic.
- \Rightarrow unstructured messages.
- Requirements:
 - ► No memory allocation except in initialization.
 - But don't use more memory than deterministic 1-halo method does.
 - Message size roughly proportional to number of departure points (up to a little metadata).
- Preferences during time stepping (not initialization):
 - Two passes to set up message structure rather than one pass plus a sort.
 - O(1) accesses with *no* hash table, just linear arrays.
 - ► No holes in arrays, i.e., unused index space.
- Strategy:
 - ► At initialization, do everything necessary to make the above feasible during time stepping.
 - ► Assemble and parse messages using these initialization-phase metadata.

```
< send/recv departure points
xs: (#x-tqt-rank
                   integer
     pad
                   i
     (lid-src-rank i
                         packed only when #x in lid > 0
      #x-in-lid i
                         > 0
      (lev
                   i
                         packed only when #x in (lid, lev) > 0
       #x
                   i
                         > 0
                 3 real
       х
        *#x) *#lev) *#lid) *#rank
qs: (q-extrema
                  2 qsize r
                               (min, max) packed together
                    qsize r
     q
      *#x) *#lev *#lid *#rank
                                  < send/recv q data
```

/* xs: (#x-tqt-rank int pad i (lid-src-rank i only packed if #x in lid > 0 #x-in-lid i > 0 only packed if #x in (lid, lev) > 0 (lev i #x i. > 0 3 real x *#x) *#lev) *#lid) *#rank */

setup_irecv(cm); // Set up to receive departure point requests.

// Determine where my departure points are. Set up requests to // remotes as well as to myself to fulfill these. Fill metadata. In parallel over my elements, determine element containing 11 departure point. Count: 11 # points in an element, 11 # points in an (element, level). 11 analyze dep points(cm, nets, nete, dep points); In parallel over remote ranks, calculate offsets and fill in 11 // xs metadata. pack dep points sendbuf pass1(cm); In parallel over elements having remotes in halo, fill x 11 // bulk data and record some offsets. pack dep points sendbuf pass2(cm, dep points);

isend(cm); // Send requests.



// While waiting, compute q extrema in each of my elements. calc_q_extrema<np>(cm, nets, nete);

// Wait for the departure point requests.
recv_and_wait_on_send(cm);



// Compute the requested q for departure points from remotes.
calc_rmt_q<np> (cm);

// Send q data, skipping messages to ranks who made 0 requests. isend(cm, false /* want_req */, true /* skip_if_empty */);



// Set up to receive q for each of my departure point requests
// sent to remotes.
setup_irecv(cm, true /* skip_if_empty */);

// While waiting to get my data from remotes, compute q for // departure points that have remained in my elements. calc_own_q<np>(cm, nets, nete, dep_points, q_min, q_max);

// Receive remote q data.
recv(cm, true /* skip_if_empty */);
// Copy these data into my data structures.
copy_q(cm, nets, q_min, q_max);

Fidelity study⁴





- Nondivergent flow test case.
- Compare (1) tuned parameters and (2) operational parameters, as in previous slide.
- SL transport is uniformly more accurate.
- For climate results, see Nov 2018 DOE Modeling PI Meeting poster³.

 $^{^{3}} https://acme-climate.atlassian.net/wiki/spaces/CNCL/pages/840073634/E8.1+Semi-Lagrangian+tracer+transport+in+the+E3SM+atmospheric+dycore$

⁴ "HOMME tuned" data are from O. Guba, et al, Optimization-based limiters for the spectral element method, JCP 2014. "CAM operational" data are from P. H. Lauritzen, et al. "Geoscientific Model Development A standard test case suite for two-dimensional linear transport on the sphere: results from a collection of state-of-the-art schemes." GMD 7(1) 2013.

Resolution: DCMIP2016 Baroclinic Instability



- Configuration: theta-1, nonhydrostatic mode, moist, ne = 30, tstep = 300, rsplit×qsplit = 6
- Eulerian at left; SL at right



(a) q_v , level 20, day 30

(b) qv, level 30, day 29



- (c) Toy chemistry tracer, level 30, day 30
- (d) Toy chemistry diagnostic, level 30, day 15